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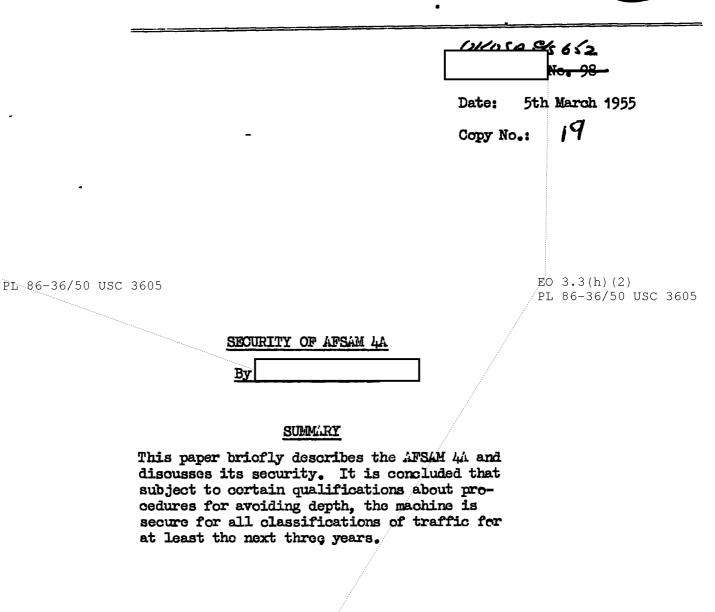
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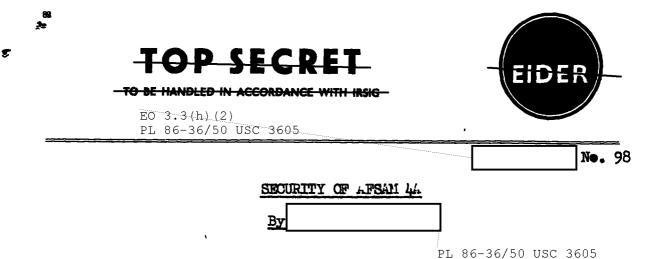
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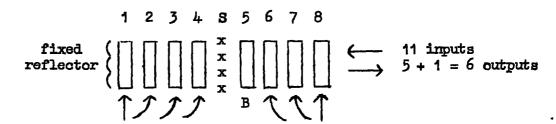
1. Introduction

AFSAM 44 is an on- or off-line additive-producing twice-through rotor machine which has recently been introduced for use in U.S. D/F and administrative networks. It is a conversion from a once-through version call M-294 or SIGNIN (BALDER and PANDORA cryptosystems). The machine and the associated cryptosystems (CENTAUR and IXION) are described and assessed in detail in BRUSA C/S 306 dated October 19th, 1953. It is there concluded that the machine, when used with these procedures, is secure for all classifications of traffic for at least the next 5 years. This paper

- (a) briefly describes the machine and the associated procedures,
- (h) discusses its security.

2. Brief Description

(a) The machine



There are 8 26-point non-reversible rotors selected from a 10-rotor set. Each rotor has a fixed alphabet and an identical fixed 5-notch notchring. The rotors are pluggable but will not be replugged in the field. R1234 and R876 form cyclonetric cascades with R1 and 8 the fast rotors. R5 is a bump rotor driven by one of the 6 outputs as in ASAM 2-1. S, between R4 and 5 is a permanent scramble-wired separator. The reflector is wired with 10 pairs and 2 triplets, as in ASAF 2-1, and the key-producing arrangements are likewise the same - 11 live inputs and 5 key outputs. The key is added level by level to 5-unit plain text, and is never transmitted in clear. The flatness of the key is discussed in Note Λ , where it is also shown that there is a slight correlation between the key and the stepping of the bump rotor.

- (b) CENTAUR
 - (i) Each pad consists of 120 pages which are assigned in blocks to stations on a net. Each page contains: (A) 4 rotor arrangements each of which has 20 tetranome designators from a sories which runs in order from 0001 to 9600 throughout the pad. (B) 176 random 4-letter groups, each designated by a digraph; the digraphs range from AA to GT on each page and are arranged alphabetically.



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(a) tł	n indicator consists of a tetranome giving nd 2 digraphs giving the alignment (as sho ne tetranome). Both tetranomes and digrap se.	the roter arrangement, wn on the same page as		
(i	Mo ar tl	or off-line use any unused tetraneme can be essages are limited to 6,000 characters. crangement can be used only once, but up t proughout the 12-hour cryptoperiod may be ent. A restart may consist either of a ne counter reset as in (iv) below.	For on-line use a rotor o 20 necessary restarts made on the same arrange-		
(iv) Co	ounter resot.			
	(,	A) Half duplex. A receive operator can c he is getting garble. The whole net g send operator steps on a randon number than 20, and transmits his character c	fes over to plain. The of positions, not less		
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	(v) Th	nere is no letter check procedure.			
(c)	IXION				
	IXION is like CENTAUR except that (i) there are about 200 pages to a pad and (ii) the tetranomes are replaced by trigraphs which are arranged alphabetically only within a single rotor arrangement.				
		a is the normal procedure. IXION will on tions warrant it.	ly be used when T/A con-		
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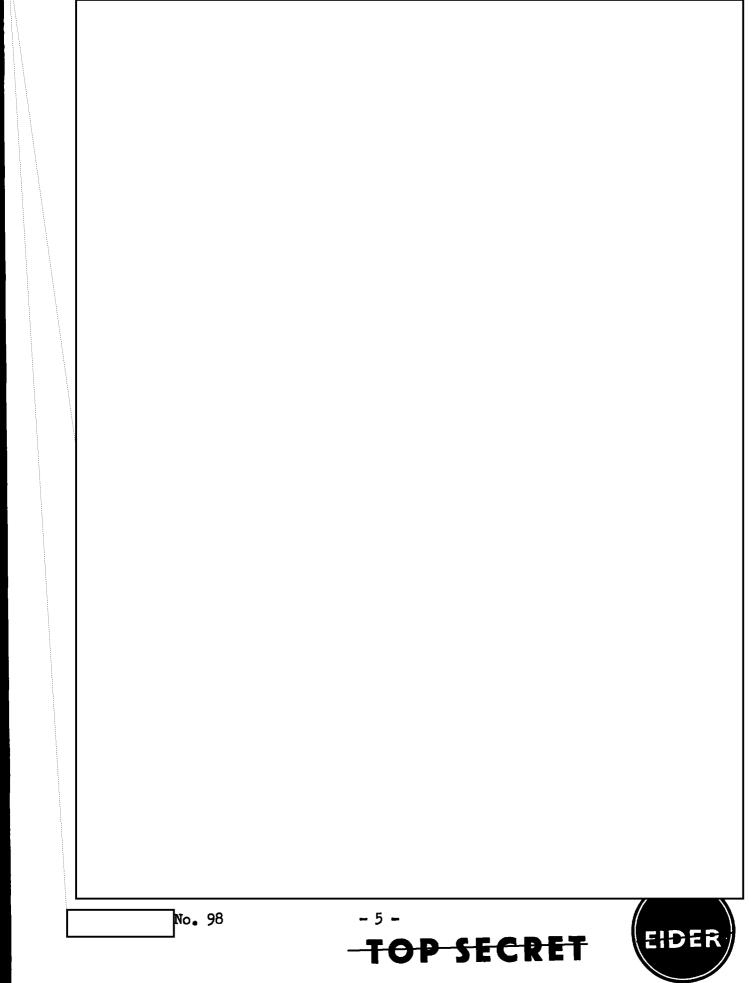
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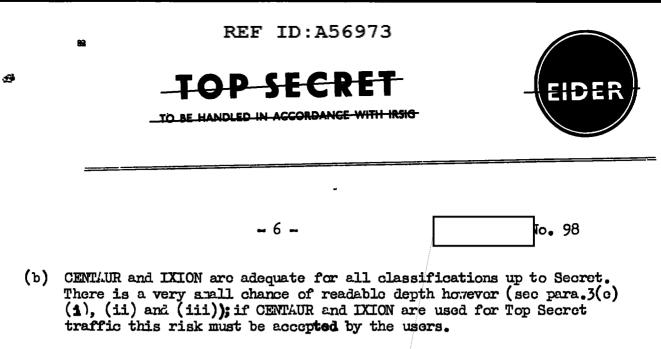
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(c) The procedure should be examined to ensure that there is no danger of depth as a result of $\mathfrak{Z}(d)(\mathbf{v})$ (B) and (C) above.

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NOTE A

Key properties

- 5. (a) In each stream the probability of cross is
 - $\frac{20}{26} \cdot \frac{11}{25} + \frac{6}{26} \begin{pmatrix} 1 \frac{14 \cdot 13}{25 \cdot 24} \end{pmatrix} = \frac{649}{1,300}$
 - (b) Counted charactervise the key is a bit rougher than this, but not seriously so. The calculation is lengthy and only the results are quoted here.

Number of orosses in key charactur	Number of characters	Probability	1 probability	P.B.
0 1	1 5	•0324 •0315	30 . 84 31.75	•0378 •0080
2	10	. 0311	32.14	0042
3	10	•0311	32.15	0043
4	5	. 0313	31.93	.0022
5	1	•0314	31.85	.0043

(c) All-dot is thus the most popular character. However approximately 86,800 key characters are needed before it can be expected to exceed its expected frequency in random material by 2 standard deviations.

6. Burnp rotor stepping

The probabilities have only been calculated for the two extreme cases:

Key	Probability that buip rotor steps
All dot	. 490
All cross	•499

A correlation of this order of size has no effect upon the security of the machine. In the case of the other 30 key characters the correlation will be smaller still.

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7. Key Sample

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The results of a count of 20,000 characters of ASAM 2-1 key are attached. Although the figures are consistent with the theoretical ones the sample size is too shall to give a significantly close correlation.

	-	
• • • • •	655	•••• x 632
x	639	xx 640
• X • • •	562	• x • • x 652
x x	668	xx. x 630
X	629	• • x • x 590
x . x	625	x . x . x 686
• x x • •	636	• x x • x 636
x x x	618	xxx, x 614
• • • X •	616	•••xx 582
x x .	607	x • • x x 599
• x • x •	592	• x • x x 604
x x . x .	643	xx.xx 684
• • x x •	622	• • x x x 600
х.хх.	598	ж _е жжж 640
• x x x •	577	• x x x x 685
x x x x .	590	xxxx 649
		20,000
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