

Extensible Security For X: Motivation and Design

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Summary

- Working towards an open source, trusted desktop.
- Need to have infrastructure for doing fine-grained access control in the X server.
- Hooks only – no specific policies.
- Local to server – no protocol changes.
- Branch development model.

What Is SELinux?

- Fine-grained Mandatory Access Control for Linux.
- Policy system based on Flask architecture.
 - Strong separation of security domains and roles.
 - Controls over process execution & resource access.
 - Diminish severity of program vulnerabilities.
- Kernel module; uses LSM security hooks.
- Some userspace changes.

SELinux Timeline

1985 LOCK (early Type Enforcement)

1990
DTMach / DTOS

1995
Utah Fluke / Flask

1999 2.2 Linux Kernel (patch)

2000

2001 2.4 Linux Kernel (patch)

2002 LSM

2003 2.6 Linux Kernel (mainline)

Present

SELinux Precursors

- LOCK
 - Early type enforcement.
- Distributed Trusted Mach (DTMach)
- Distributed Trusted OS (DTOS)
 - Improved design and implementation in Mach.
- Flux Advanced Security Kernel (Flask)
 - Flexible MAC architecture in the Flux OS.

SELinux Distributions

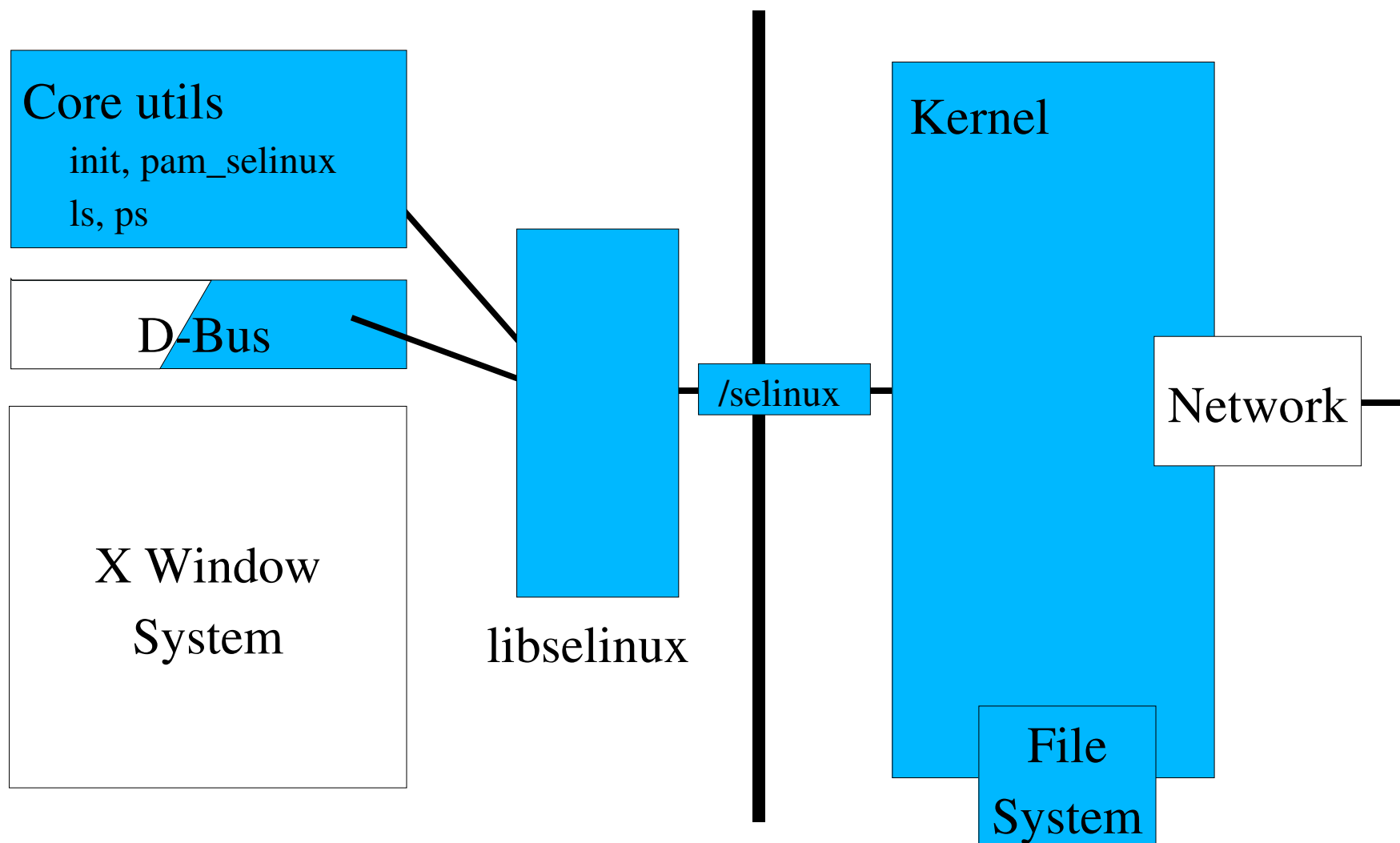
- Fedora Core 2
- Hardened Gentoo
- Debian (packages)
- SE-BSD (port)
- SE-Darwin (port)

SELinux Research Agenda

- Security architecture research
- Kernel prototype code
- Kernel production code
- Userspace enhancements
 - Local GUI security
- Labeled networking
- Network-wide policy



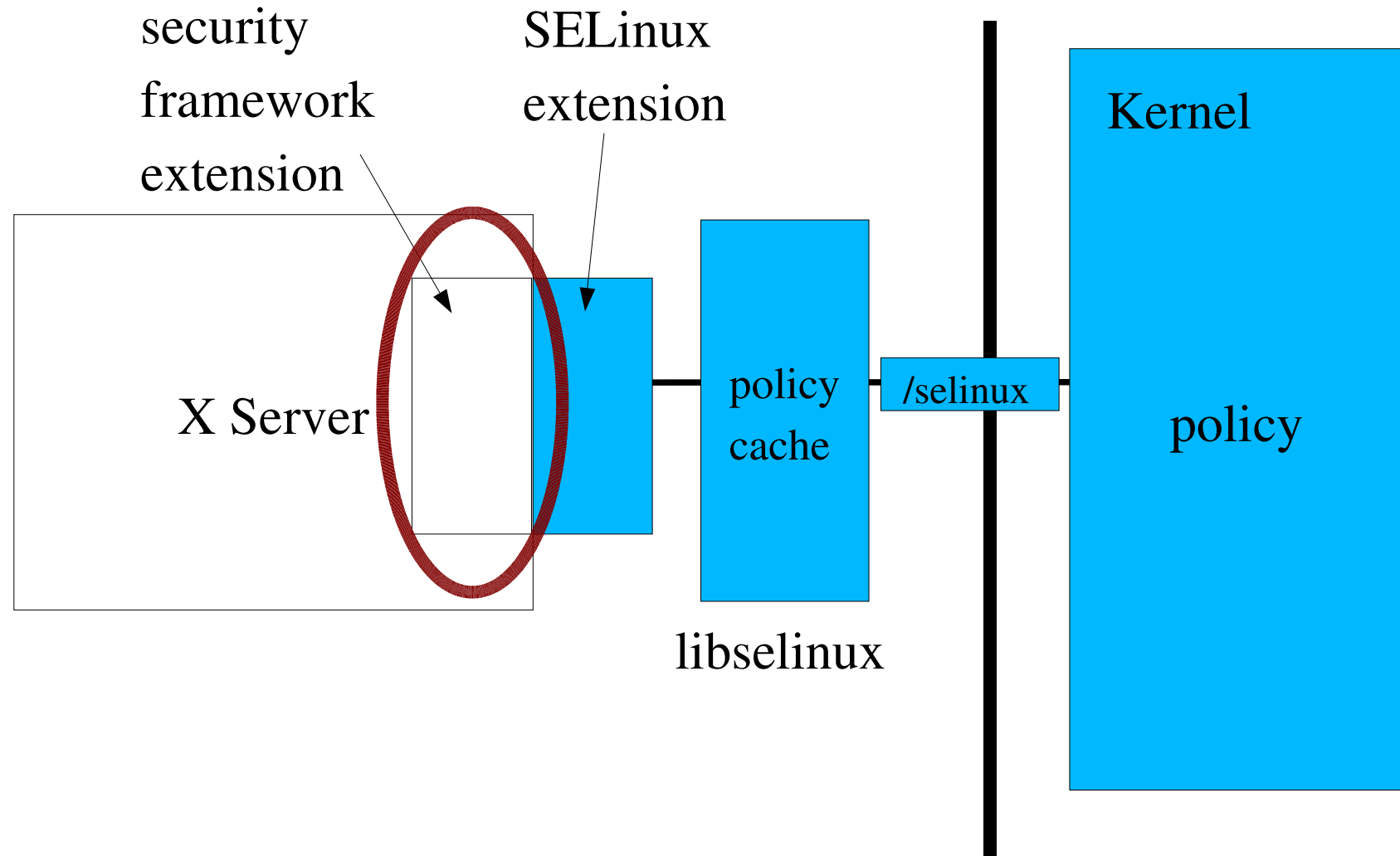
Current State of SELinux



GUI Security

- GUI security is the last piece of the complete SELinux desktop system.
- X Window System operations should be policy-controlled.
- Need to write policy for the X Window System and have the X server enforce it.
- Generalize: make it easy to write access control extensions for the X server.

SELinux/X architecture



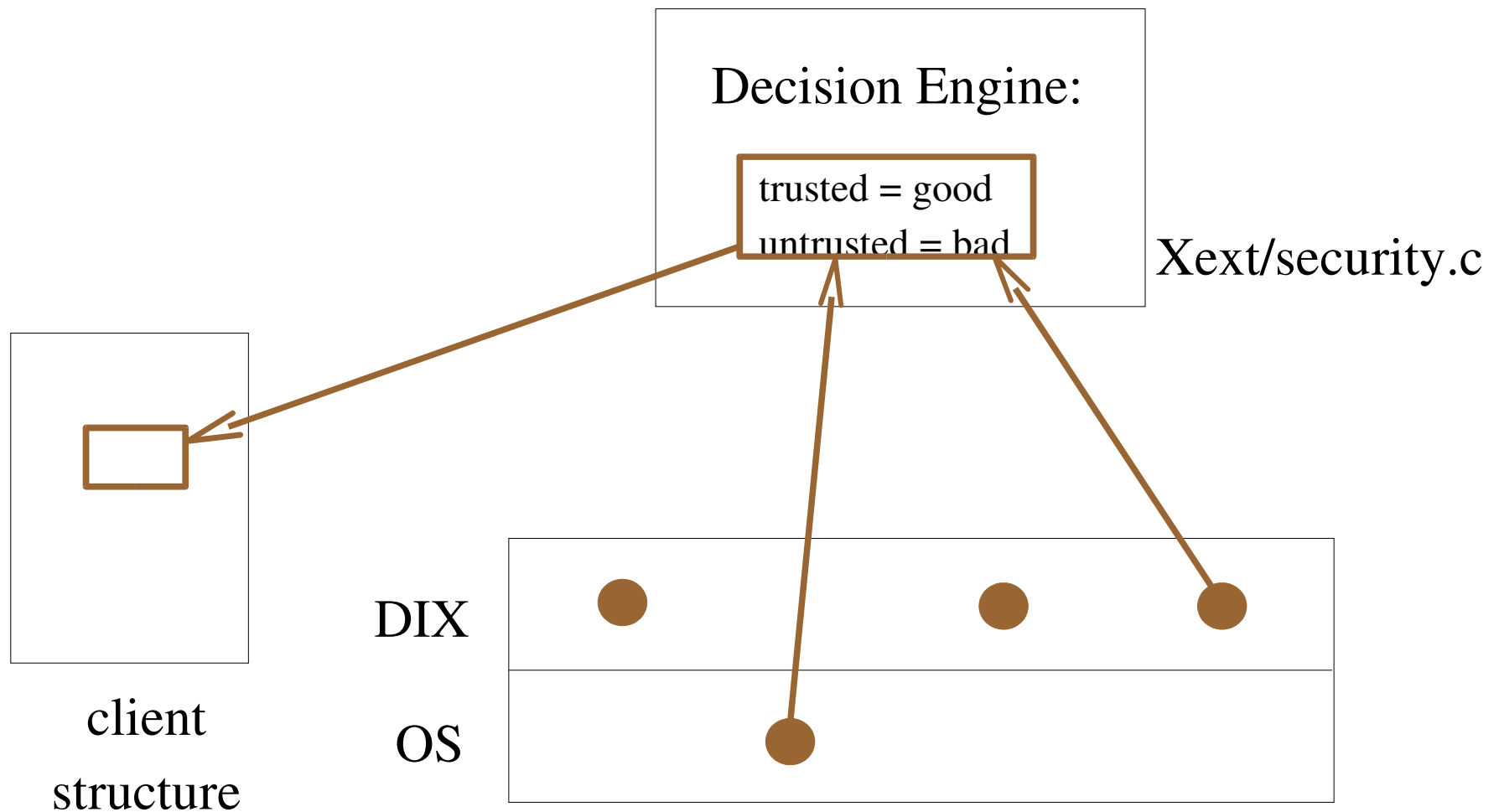
X Security and the Network

- SELinux is currently a local system.
- SELinux does not have labeled networking or network authentication.
- X Window System big problem is client authentication over the network.
- Local security engine, new auth solution can be independent; complementary.

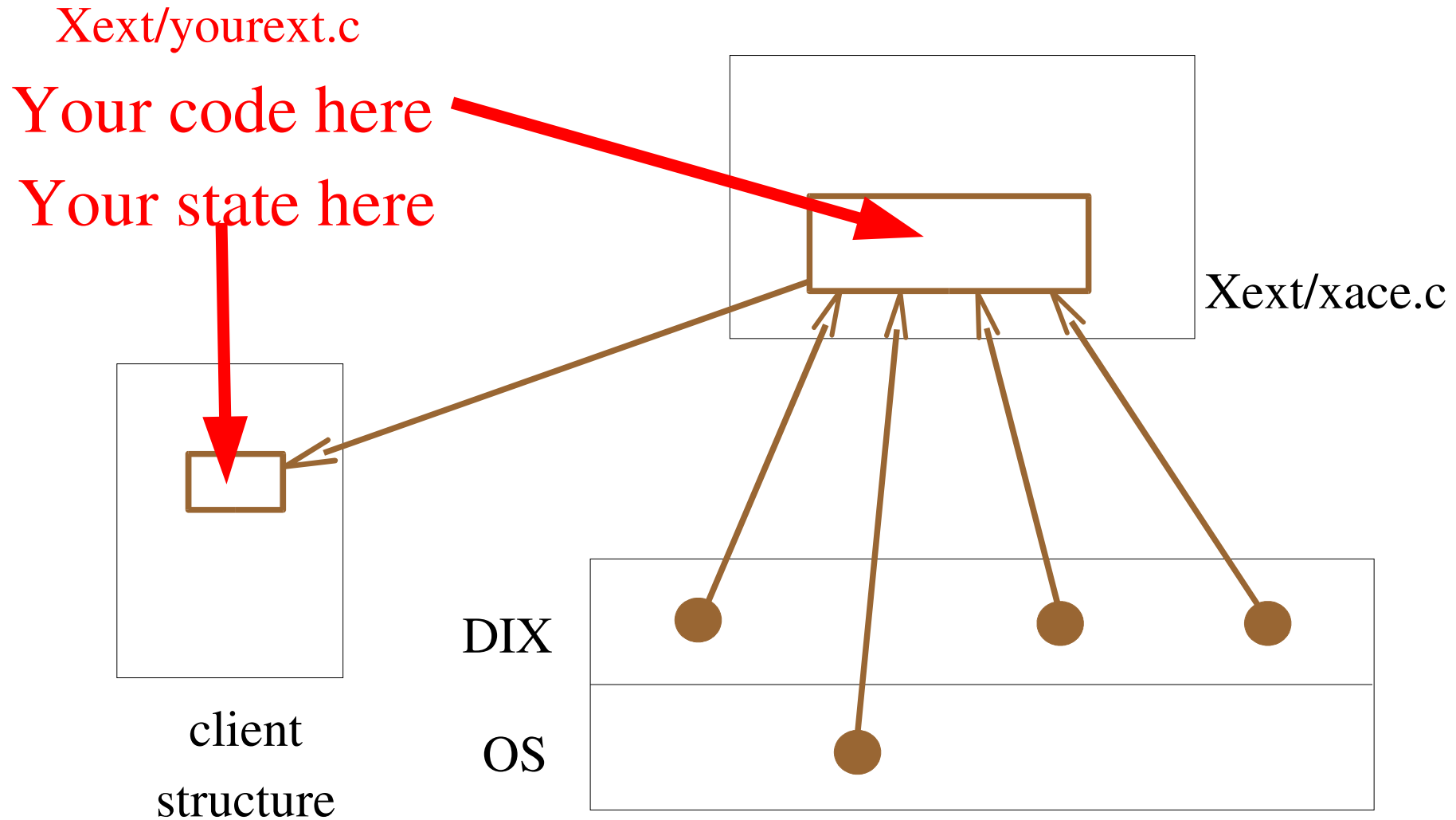
Goals for Security Framework

- Based on existing work.
- Easily extensible.
- Non-intrusive: based on callbacks, not local code.
- Works at dispatch (DIX) layer to avoid performance issues.
- Provides framework for arbitrary decision-making (access control) extensions.

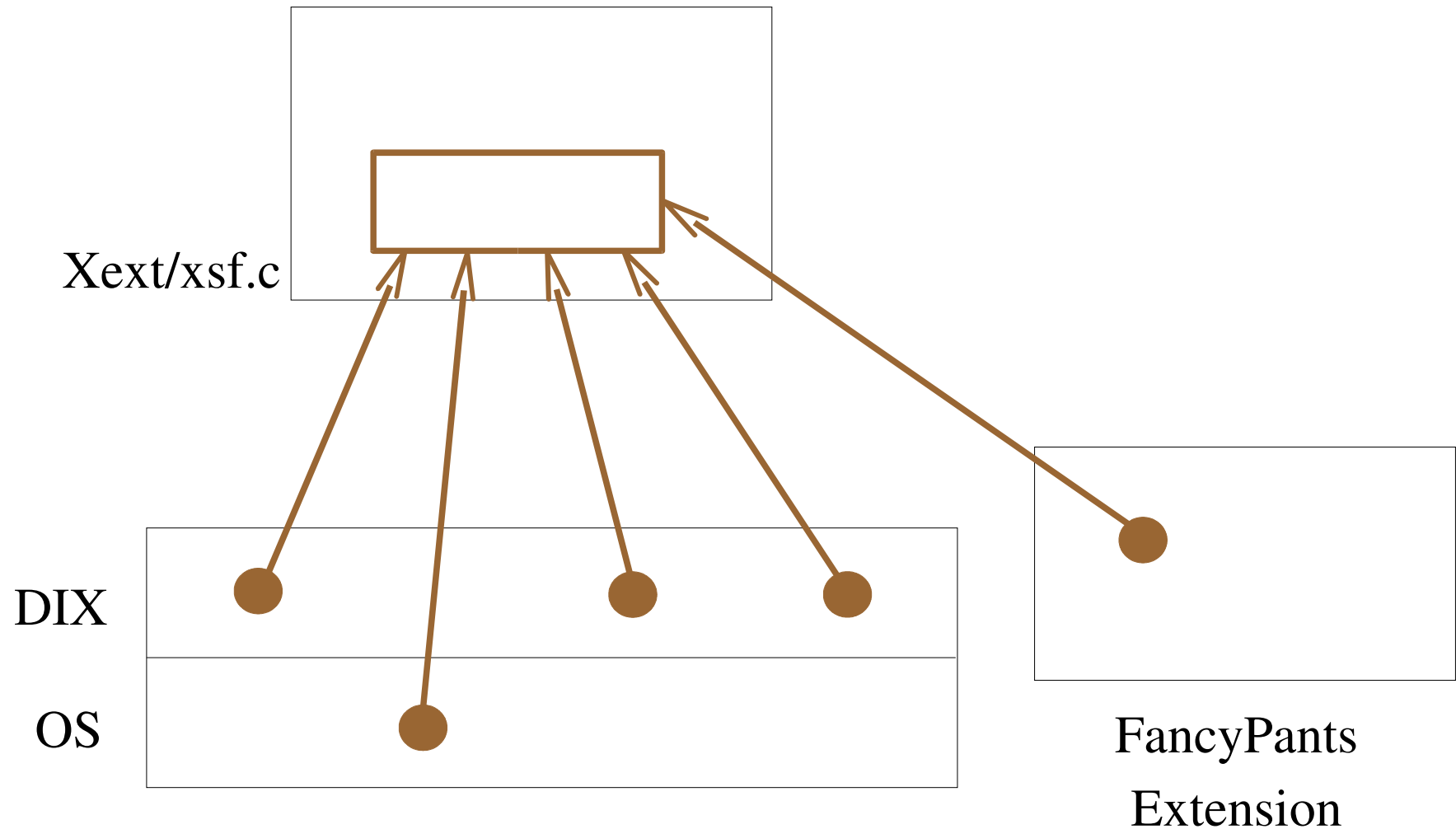
Current XC-Security Extension



Generalized Security Extension



Easily Extensible



Non-Intrusive

- At decision point, only need to pass parameters to a hook function and check the result.
- Actual security code is in the callback functions.
- Separates security code from the core code.
- Whole framework is compile-time option.

Code Examples

```
ProcDoSomething(...)
{
    rval = SecurityLookupIDByType
(client,
        MyResType, stuff->id,
        SecurityReadAccess);


---


    if (!rval) return BadSomething;
ProcDoSomething(...)
    DoNiftyStuff();
}
#ifdef XACE
    if (!SecurityHook
(XACE_FOO_ACCESS,
        client, whatever))
        return BadSomething;
```

Sample Hooks

CORE_DISPATCH

EXT_DISPATCH

- Replace XC-Security shadow dispatcher.

RESOURCE_ACCESS

DEVICE_ACCESS

PROPERTY_ACCESS

- Replace SecurityCheck*Access() functions.

MAP_ACCESS

BACKGRND_ACCESS

- Replace untrusted child & background “None” checks.

Performance Issues

- Keep hooks at the DIX layer.
- $O(1)$ hook calls per protocol request.
- Make decision before starting graphics operation.

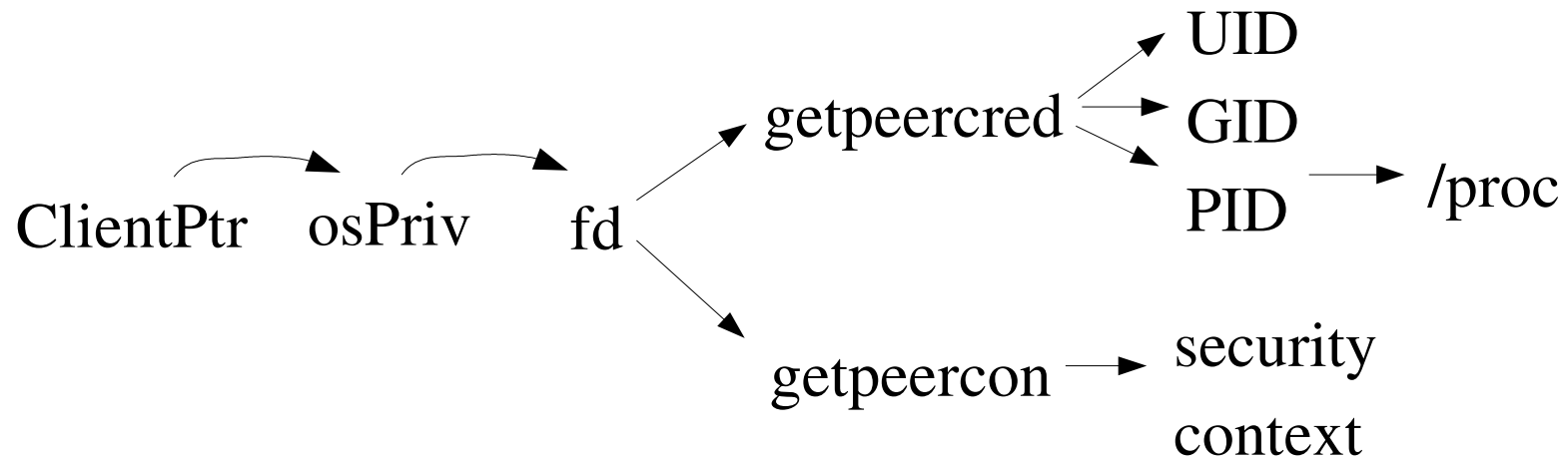
Provides General Framework

- Arbitrary new extensions can be written to use the framework's interface.
 - Provide own state for server objects and own callback functions.
- No client-side work necessary (except for proper error handling).

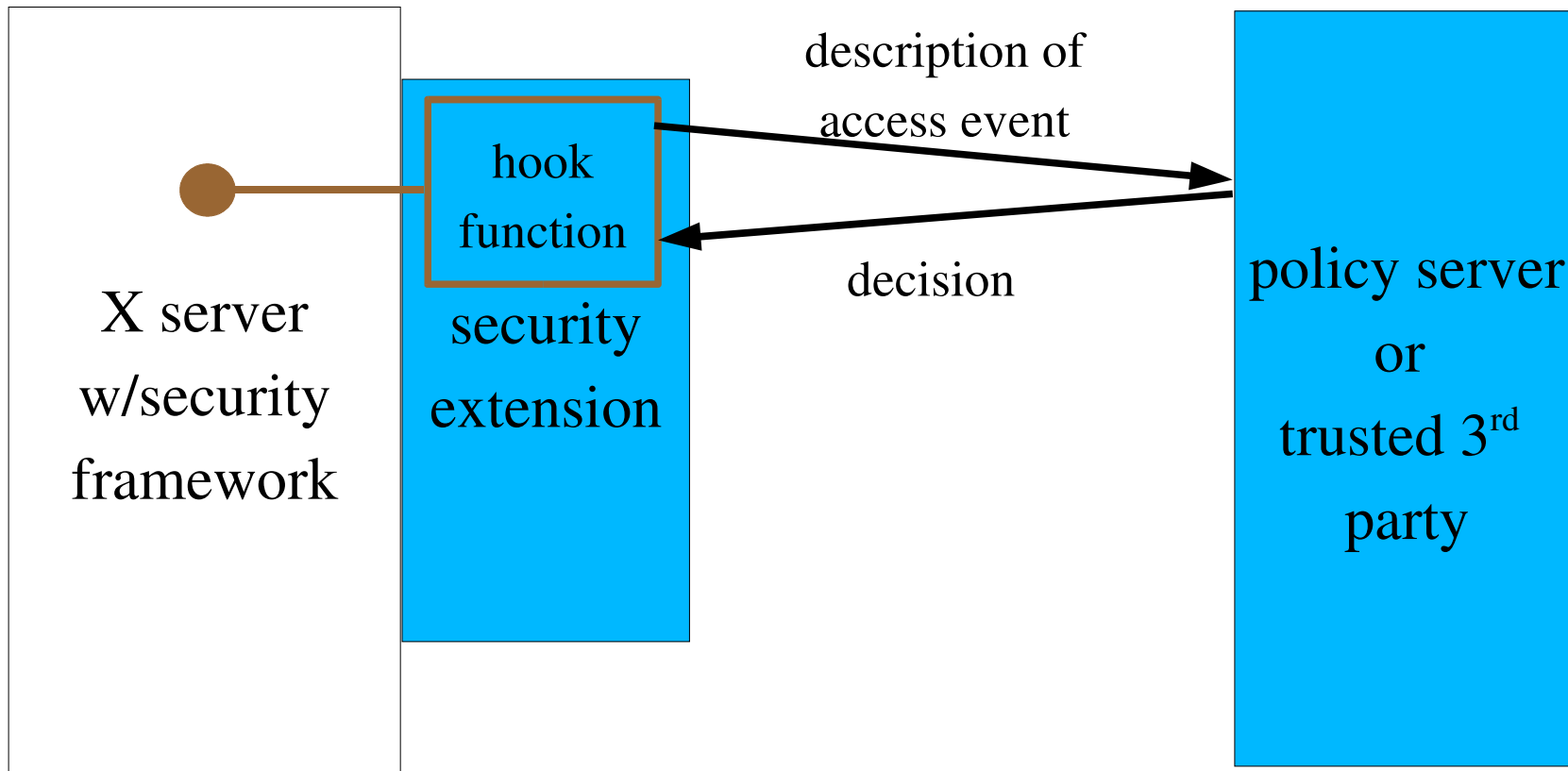
How to make Security Decisions?

- Need information about the connected client.
- Obtain once - store as client state.
- Can get:
 - From the local system.
 - From the system security policy.
 - From the authentication mechanism.

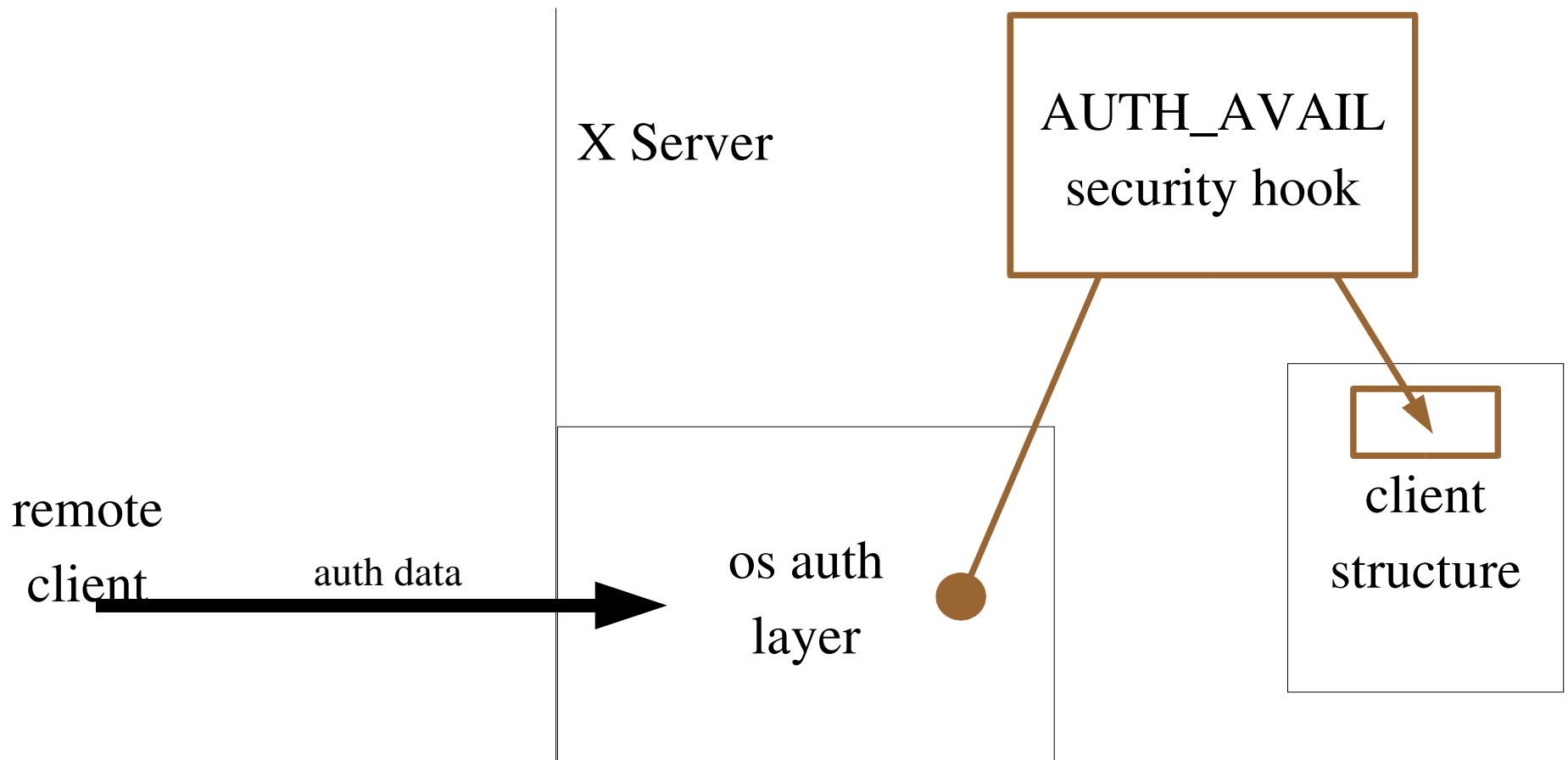
Local System



Local Security Policy



Authentication Protocol



Authentication Protocol, cont'd.

- Opportunity to combine power of the security framework with new, secure authentication methods.
- Design protocol, then write security extension to do fine-grained access control.
- At connect time, pass auth data to a security hook.
- Callbacks on that hook can set client state based on the auth data

Other Security Issues

- Trusted window labeling
 - Pass some String label to window manager on request.
 - Define a standard way to do this (new extension).
 - Or, use a Property on the window (that other clients can't mess with).

In Closing

- Flexible MAC on the open-source desktop is within reach.
- Generalized security engine, as described, will benefit SELinux project and others.
- Combine with better authentication for full solution.

Contact Information

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